

# **Lecture 7**

## **Instruction Scheduling**

- I      Basic Block Scheduling
- II     Global Scheduling (for Non-Numeric Code)

Reading: Chapter 10.3 - 10.4

# I. Scheduling Constraints

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- **Data dependences**

- The operations must generate the same results as the corresponding ones in the original program.

- **Control dependences**

- All the operations executed in the original program must be executed in the optimized program

- **Resource constraints**

- No over-subscription of resources.

# Data Dependence

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- **Must maintain order of accesses to potentially same locations**

- True dependence: write -> read (RAW hazard)

a = ...  
= a

- Output dependence: write -> write (WAW hazard)

a = ...  
a = ...

- Anti-dependence: read -> write (WAR hazard)

= a  
a =

- **Data Dependence Graph**

- Nodes: operations
- Edges:  $n_1 \rightarrow n_2$  if  $n_2$  is data dependent on  $n_1$   
labeled by the execution length of  $n_1$

# Analysis on Memory Variables

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- **Undecidable in general**

```
read x
read y
A[x] = ...
...   = A[y]
```

- **Two memory accesses can potentially be the same unless proven otherwise**

- **Classes of analysis**

- simple:  $\text{base} + \text{offset1} = \text{base} + \text{offset2}?$
- “data dependence analysis”:  
Array accesses whose indices are affine expressions of loop indices  
 $A[2i] = A[2i+1]?$
- interprocedural analysis:  $\text{global} = \text{parameter}?$
- pointer analysis:  $\text{pointer1} = \text{pointer2}?$

- **Data dependence analysis is useful for many other purposes**

# Resource Constraints

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- Each instruction type has a resource reservation table

		Functional units					
		ld	st	alu	fmpy	fadd	br ...
Time	0						
	1						
	2						

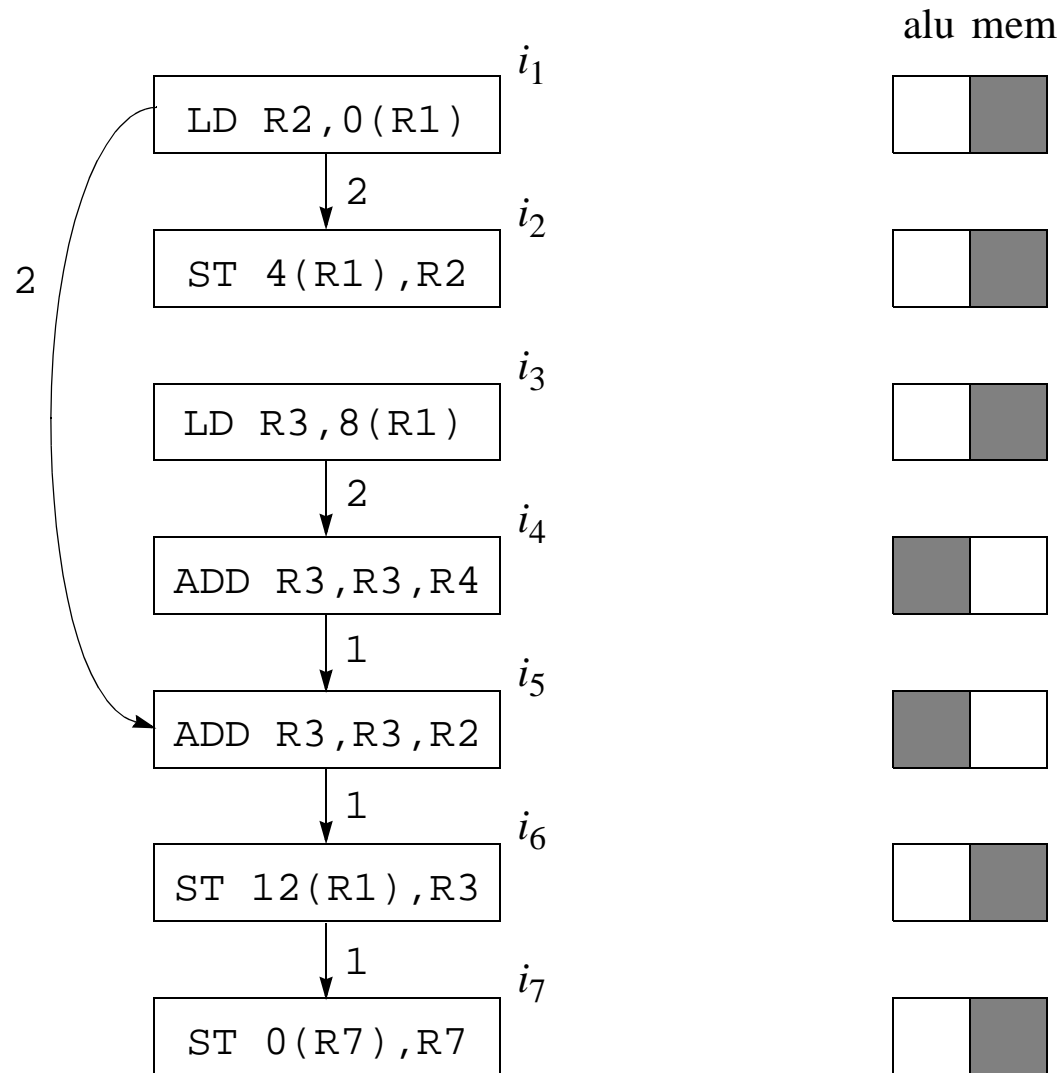
- Pipelined functional units: occupy only one slot
- Non-pipelined functional units: multiple time slots
- Instructions may use more than one resource
- Multiple units of same resource
- Limited instruction issue slots  
may also be managed like a resource

# Example of a Machine Model

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- Each instruction can execute 2 operations
- 1 ALU operation or branch operation  
`Op dst, src1, src2` executes in 1 clock
- 1 load or store operation  
`LD dst, addr` result is available in 2 clocks  
pipelined: can issue `LD` next clock  
`ST src, addr` executes in 1 clock cycle

# Basic Block Scheduling



# With Resource Constraints

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- NP-complete in general => Heuristics time!
- List Scheduling

READY = nodes with 0 predecessors

Loop until READY is empty {

Let **n** be the node in READY with **highest priority**

Schedule **n** in the earliest slot  
that satisfies precedence + resource constraints

Update predecessor count of **n**'s successor nodes  
Update READY

}



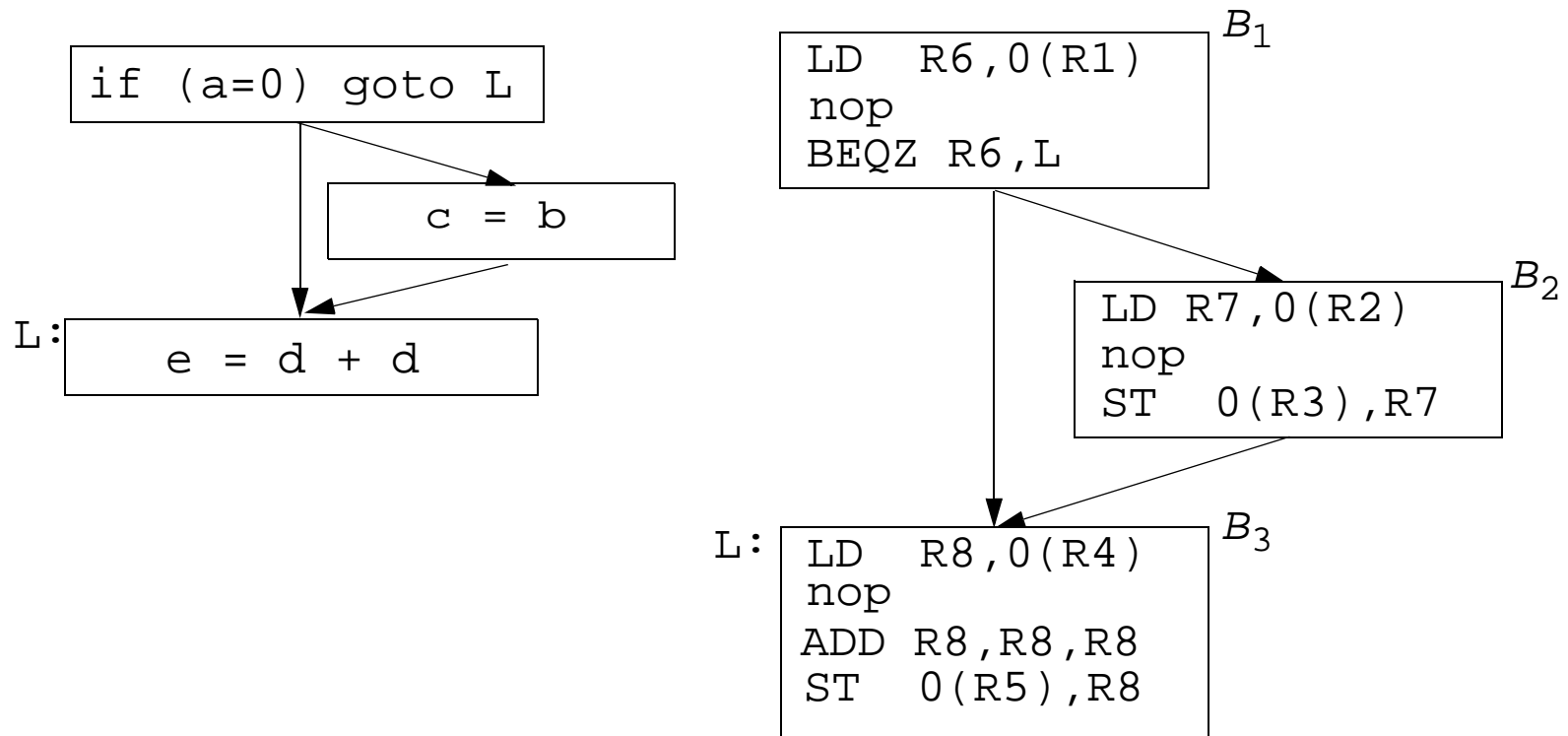
# List Scheduling

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- **Scope: DAGs**
  - Schedules operations in topological order
  - Never backtracks
- **Variations**
  - Priority function for node **n**
    - delay: max delay slots from **n** to any node
    - critical path: max clocks from **n** to any node
    - resource requirements
    - source order

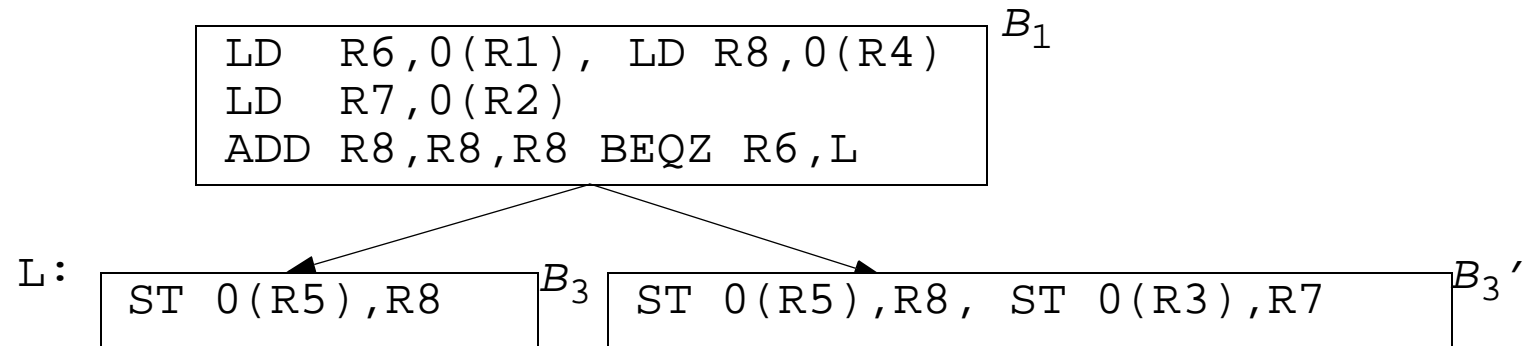
## II. Introduction to Global Scheduling

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# Result of Code Scheduling

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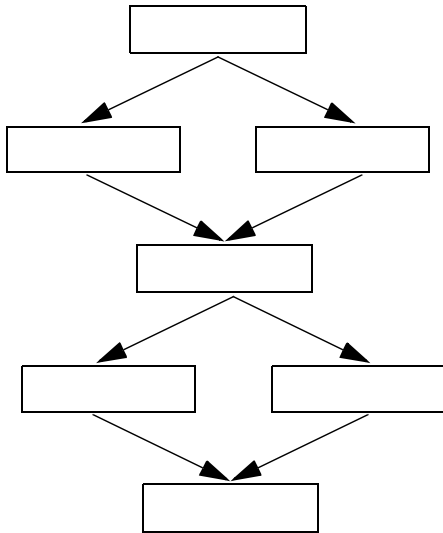


# Terminology

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## Control equivalence

- Two operations  $o_1$  and  $o_2$  are control equivalent if  $o_1$  is executed if and only if  $o_2$  is executed.



## Control dependence

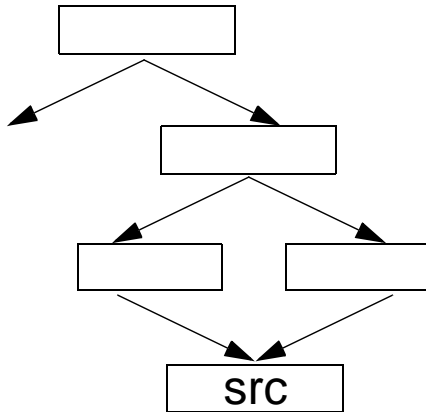
- An op  $o_2$  is control dependent on op  $o_1$  if the execution of  $o_2$  depends on the outcome of  $o_1$ .

## Speculation

- An operation  $o_1$  is speculatively executed if it is executed before all the operations it control-dependent upon have been executed.
- No exception, satisfy data dependences

# Code Motions

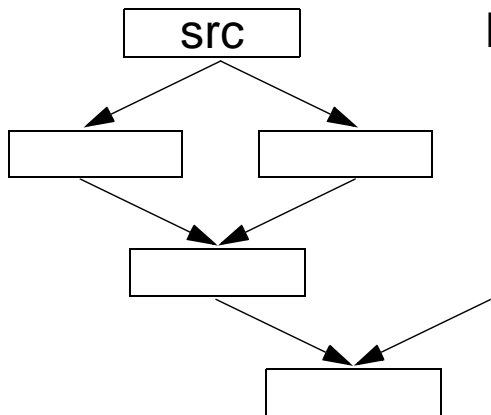
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Goal: Shorten execution time probabilistically

Moving instructions up

- Move instruction to a cut set (from entry)
- Speculation: even when not anticipated.

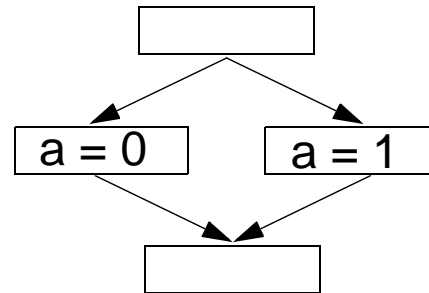


Moving instructions down

- Move instruction to a cut set (from exit)
- May execute extra instruction
- Can duplicate code

# A Note on Data Dependences

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# General-Purpose Applications

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- **Lots of data dependences**
- **Key performance factor: memory latencies**
- **Move memory fetches up**
  - Speculative memory fetches can be expensive
- **Control-intensive: get execution profile**
  - Static estimation
    - Innermost loops are frequently executed: back edges are likely to be taken
    - Edges that branch to exit and exception routines are not likely to be taken
  - Dynamic profiling
    - Instrument code and measure using representative data

# A Basic Global Scheduling Algorithm

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- **Schedule innermost loops first**
- **Only upward code motion**
- **No creation of copies**
- **Only one level of speculation**



# Program Representation

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- **A region in a control flow graph is**
  - a set of basic blocks and all the edges connecting these blocks,
  - such that control from outside the region must enter through a single entry block.
- **A function is represented as a hierarchy of regions**
  - The whole control flow graph is a region
  - Each natural loop in the flow graph is a region
  - Natural loops are hierarchically nested
- **Schedule regions from inner to outer**
  - treat inner loop as a black box unit,  
can schedule around it but not into it
  - ignore all the loop back edges --> get an acyclic graph

# Algorithm

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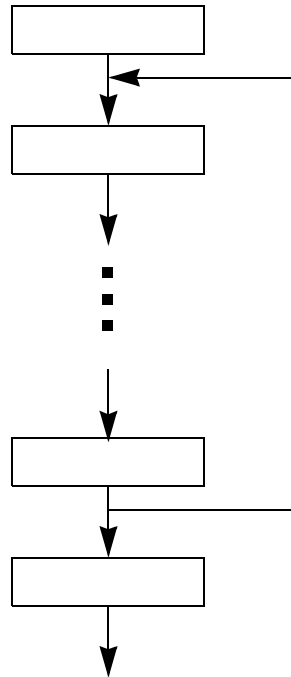
```
Compute data dependences;
For each region from inner to outer {
  For each basic block B in prioritized topological order {
    CandBlocks = ControlEquiv{B}  $\cup$ 
                  Dominated-Successors{ControlEquiv{B}};
    CandInsts = ready operations in CandBlocks;
    For (t = 0, 1, ... until all operations from B are scheduled {
      For (n in CandInst in priority order) {
        if (n has no resource conflicts at time t) {
          S(n) = < B, t >
          Update resource commitments
          Update data dependences
        }
      }
      Update CandInsts;
    }
  }
}
```

- **Priority functions**
  - Non-speculative before speculative

# Extensions

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- **Prepass before scheduling: loop unrolling**
- **Especially important to move operation up loop back edges**



# Summary

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- **List scheduling**
- **Global scheduling**
  - Legal code motions
  - Heuristics