

Lecture 7

Instruction Scheduling

- I Basic Block Scheduling
- II Global Scheduling (for Non-Numeric Code)

Reading: Chapter 10.3 - 10.4

I. Scheduling Constraints

- **Data dependences**
 - The operations must generate the same results as the corresponding ones in the original program.
- **Control dependences**
 - All the operations executed in the original program must be executed in the optimized program
- **Resource constraints**
 - No over-subscription of resources.

Data Dependence

- **Must maintain order of accesses to potentially same locations**

- True dependence: write -> read (RAW hazard)

```
a = ...  
= a
```

- Output dependence: write -> write (WAW hazard)

```
a = ...  
a = ...
```

- Anti-dependence: read -> write (WAR hazard)

```
= a  
a =
```

- **Data Dependence Graph**

- Nodes: operations
- Edges: $n_1 \rightarrow n_2$ if n_2 is data dependent on n_1
labeled by the execution length of n_1

Analysis on Memory Variables

- **Undecidable in general**

```
read x  
read y  
A[x] = ...  
... = A[y]
```

- **Two memory accesses can potentially be the same unless proven otherwise**

- **Classes of analysis**

- simple: $\text{base+offset1} = \text{base+offset2?}$
- “data dependence analysis”:
Array accesses whose indices are affine expressions of loop indices
 $A[2i] = A[2i+1]?$
- interprocedural analysis: $\text{global} = \text{parameter?}$
- pointer analysis: $\text{pointer1} = \text{pointer2?}$

- **Data dependence analysis is useful for many other purposes**

Resource Constraints

- Each instruction type has a resource reservation table

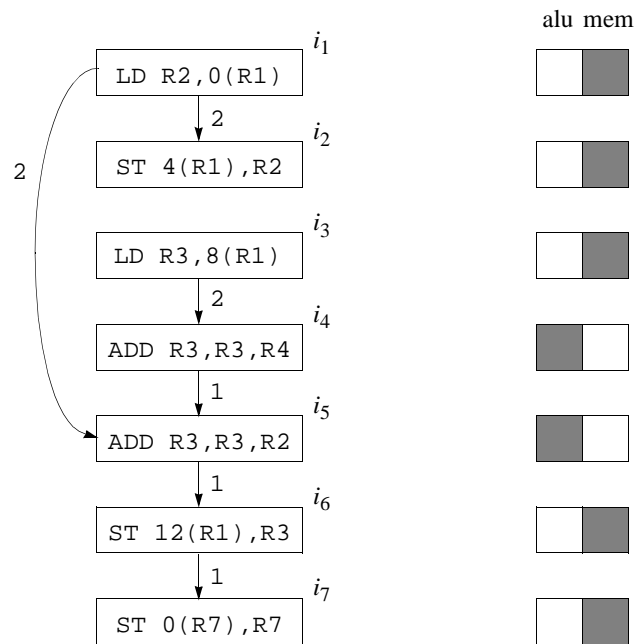
		Functional units					
		ld	st	alu	fmpyfadd	br	...
Time	0						
	1						
	2						

- Pipelined functional units: occupy only one slot
- Non-pipelined functional units: multiple time slots
- Instructions may use more than one resource
- Multiple units of same resource
- Limited instruction issue slots
may also be managed like a resource

Example of a Machine Model

- Each instruction can execute 2 operations
- 1 ALU operation or branch operation
`Op dst, src1, src2` executes in 1 clock
- 1 load or store operation
`LD dst, addr` result is available in 2 clocks
pipelined: can issue LD next clock
`ST src, addr` executes in 1 clock cycle

Basic Block Scheduling



With Resource Constraints

- NP-complete in general => Heuristics time!
- List Scheduling

READY = nodes with 0 predecessors

Loop until READY is empty {

Let n be the node in READY with **highest priority**

Schedule n in the earliest slot
that satisfies precedence + resource constraints

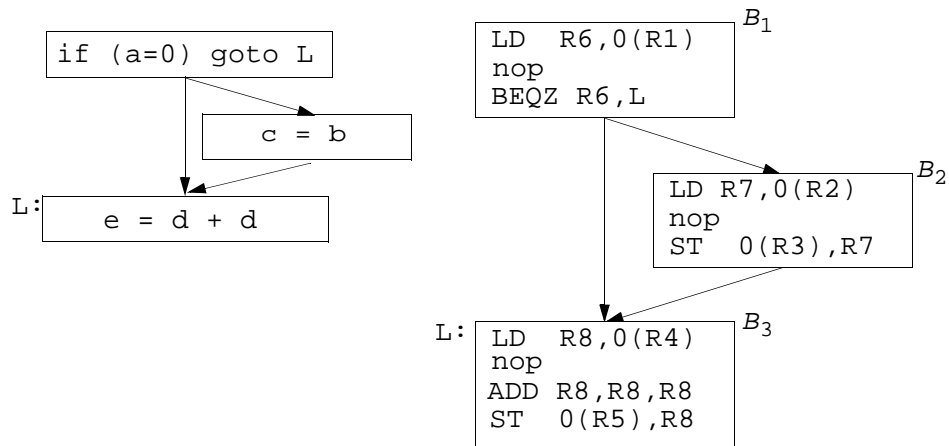
Update predecessor count of n 's successor nodes
Update READY

}

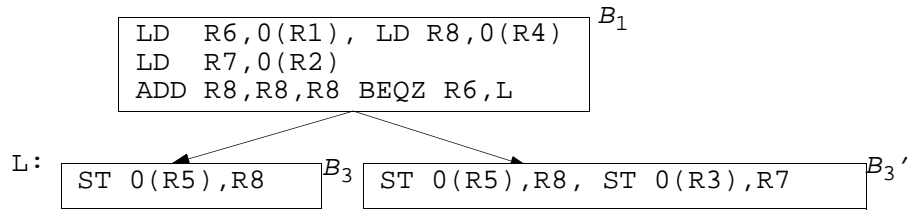
List Scheduling

- **Scope: DAGs**
 - Schedules operations in topological order
 - Never backtracks
- **Variations**
 - Priority function for node **n**
 - delay: max delay slots from **n** to any node
 - critical path: max clocks from **n** to any node
 - resource requirements
 - source order

II. Introduction to Global Scheduling



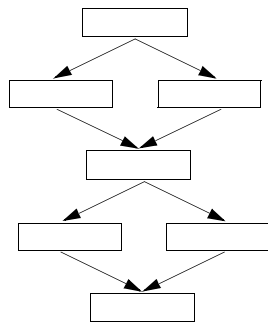
Result of Code Scheduling



Terminology

Control equivalence

- Two operations o_1 and o_2 are control equivalent if o_1 is executed if and only if o_2 is executed.



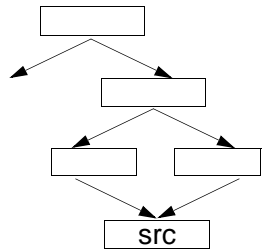
Control dependence

- An op o_2 is control dependent on op o_1 if the execution of o_2 depends on the outcome of o_1 .

Speculation

- An operation o_1 is speculatively executed if it is executed before all the operations it control-dependent upon have been executed.
- No exception, satisfy data dependences

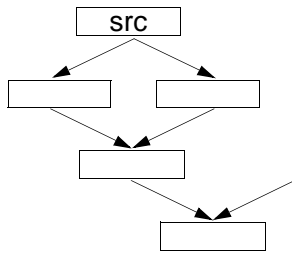
Code Motions



Goal: Shorten execution time probabilistically

Moving instructions up

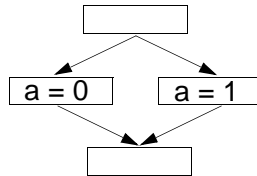
- Move instruction to a cut set (from entry)
- Speculation: even when not anticipated.



Moving instructions down

- Move instruction to a cut set (from exit)
- May execute extra instruction
- Can duplicate code

A Note on Data Dependences



General-Purpose Applications

- **Lots of data dependences**
- **Key performance factor: memory latencies**
- **Move memory fetches up**
 - Speculative memory fetches can be expensive
- **Control-intensive: get execution profile**
 - Static estimation
 - Innermost loops are frequently executed: back edges are likely to be taken
 - Edges that branch to exit and exception routines are not likely to be taken
 - Dynamic profiling
 - Instrument code and measure using representative data

A Basic Global Scheduling Algorithm

- **Schedule innermost loops first**
- **Only upward code motion**
- **No creation of copies**
- **Only one level of speculation**

Program Representation

- **A region in a control flow graph is**
 - a set of basic blocks and all the edges connecting these blocks,
 - such that control from outside the region must enter through a single entry block.
- **A function is represented as a hierarchy of regions**
 - The whole control flow graph is a region
 - Each natural loop in the flow graph is a region
 - Natural loops are hierarchically nested
- **Schedule regions from inner to outer**
 - treat inner loop as a black box unit, can schedule around it but not into it
 - ignore all the loop back edges --> get an acyclic graph

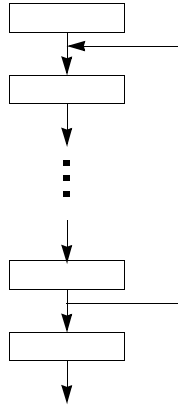
Algorithm

```
Compute data dependences;
For each region from inner to outer {
  For each basic block B in prioritized topological order {
    CandBlocks = ControlEquiv{B} ∪
                Dominated-Successors{ControlEquiv{B}};
    CandInsts = ready operations in CandBlocks;
    For (t = 0, 1, ... until all operations from B are scheduled {
      For (n in CandInst in priority order) {
        if (n has no resource conflicts at time t) {
          S(n) = < B, t >
          Update resource commitments
          Update data dependences
        }
      }
    }
    Update CandInsts;
  }
}
```

- **Priority functions**
 - Non-speculative before speculative

Extensions

- **Prepass before scheduling: loop unrolling**
- **Especially important to move operation up loop back edges**



Summary

- **List scheduling**
- **Global scheduling**
 - Legal code motions
 - Heuristics