# joeq Compiler System

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## Plan for Today

- Joeq System Overview
- Lifecycle of Analyzed Code
- Source Code Representation
- Writing and Running a Pass
- Assignment: Dataflow Framework

#### Background on joeq

- □ A compiler system for analyzing Java code
  - Developed by John Whaley and others
  - Used on a daily basis by the SUIF compiler group
  - An infrastructure for many research projects: 10+ papers rely on joeq implementations
- □ Visit http://joeq.sourceforge.net for more
- □ Or read

http://www.stanford.edu/~jwhaley/papers/ivme03.pdf

#### joeq Design Choices

- □ Most of the system is implemented in pure Java
- □ Thus, analysis framework and bytecode processors work everywhere
- □ We treat joeq as a front-end and middle-end
- ☐ But it can be used as a VM as well
  - System-specific code is patched in when the joeq system compiles itself or its own runtime
  - These are ordinary C routines
  - Systems supported by full version: Linux and Windows under x86

#### joeq Components

Full system is very large:	Synchronization
~100,000 lines of code	Assembler
Allocator	Class Library
Bootstrapper	Compiler (Bytecode)
Classfile structure	Debugger
Compiler (Quad)	Bytecode Interpreters
Garbage Collector	Linkers
Quad Interpreters	Reflection support
Memory Access	Scheduling
Safe/Unsafe barriers	UTF-8 Support

We restrict ourselves to only the compiler and classfile routines, which is closer to 40,000 lines of code

# Starting at the Source

#### Lifecycle of Analyzed Code

- □ Everything begins as source code
- □ A very "rich" representation
  - Good for reading
  - Hard to analyze
- □ Lots of high-level concepts here with (probably) no counterparts in the hardware
  - Virtual function calls
  - Direct use of monitors and condition variables
  - Exceptions
  - Reflection
  - Anonymous classes
  - Threads

#### Source to Bytecode

- □ javac or jikes compiles source into a machine-independent bytecode format
- □ Coarse structure of the program is still maintained
  - Each class is a file
  - Split up into methods and fields
  - The bytecodes themselves are stored as a member attribute in methods that have them
  - Bytecode instructions are themselves high level:
    - □ invokevirtual
    - monitorenter
    - □ arraylength

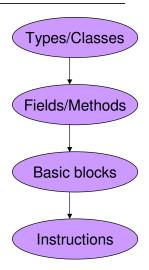
#### Analysis and Source Code

- □ No need to bother with source code, since structure is maintained in classfile format
- □ Moreover, bytecode is indifferent to language changes
- □ Reading in code:
  - 1. joeq finds and loads requested files through the classpath
  - 2. Each source component in the classfile has a corresponding object:
    - □ jq\_Class
    - □ jq\_Method
    - □ etc.
- Method bodies are transformed from bytecode arrays to more convenient representation

# How Source Code is Represented Within joeq

#### Source Code Representation

- □ joeq is designed primarily to work with Java
  - Operates at all levels of abstraction
  - Has classes corresponding to each language component
- Relevant packages in joeq
  - joeq.Class package: classes that represent Java source components (classes, fields, methods, etc.)
  - joeq.Compiler.BytecodeAnalysis package: analysis of Java bytecode
  - joeq.Compiler.Quad package: Classes relevant to joeq's internal "quad" format
- □ Be careful with your imports
  - avoid name conflicts with java.lang.Class and java.lang.Compiler



#### joeq.Class: Types and Classes

- □ jq\_Type: corresponds to any Java type
- □ jq\_Primitive: static elements representing primitive types
- □ jq\_Array: multidimensional arrays that have a component type, which itself is a jq\_Array
- □ jq\_Class: a defined class

#### joeq.Class: Fields and Methods

- □ Subclasses of jq\_Field and jq\_Method
  - Class hierarchy distinguishes between instance and class (static) members, but this detail is generally hidden from higher analyses
- □ Access to the types hierarchy: declaring types, parameter/return types, etc.
- □ Names are stored as UTF.Utf8 objects, so convert with toString() to make use of them

#### Analyzing Bytecode

- ☐ The Java Virtual Machine stores program code as *bytecodes* that serve as instructions to a stack machine of sorts
- □ Raw material for all analysis of Java code
- □ Preserves vast amounts of source information:
  - De-compilers can reconstruct source almost perfectly, down to variable names and line numbers

### Example of Java Bytecode

```
class ExprTest {
                                     int test(int);
                                      Code:
    int test(int a) {
                                       0:
                                            iload_1
       int b, c, d, e, f;
                                       1:
                                            bipush 10
       c = a + 10;
                                            iadd
       f = a + c;
                                        4:
                                            istore_3
                                       5:
                                            iload_1
                                            iload_3
                                        6:
       if(f > 2){
                                       7:
                                            iadd
         f = f - c;
                                            istore 6
       }
                                       10: iload
                                       12: iconst_2
                                       13: if_icmple
                                                           22
       return f;
                                       16: iload
     }
                                       18: iload_3
}
                                        19: isub
                                       20: istore 6
    javac test.java
                                       22: iload
24: ireturn
    javap -c ExprTest
```

#### Bytecode Details

- ☐ The implied running model of the Java Virtual Machine is a stack machine
  - Local variables correspond to registers
  - Computation occurs on a stack
- ☐ This is hard to analyze!
- ☐ Fortunately, the JVM requires that bytecode pass strict type-checking and stack consistency checking
- ☐ **Gosling Property:** At each instruction, the types of every element on the stack, and every local variable, are all well defined
- □ By extension, the stack must have a specific height at each program point

#### Converting Bytecodes to Quads

- □ joeq converts bytecodes to four-address code, called "Quads"
- ☐ The highly abstract bytecode instructions have Quad counterparts
- One operator, up to four operands
   OPERATOR OP1 OP2 OP3 OP4
- □ Approximately 100 operators and 15 varieties of operands
- □ Details on the quads and relevant methods are on the course website's joeq documentation:
  - http://suif.stanford.edu/~courses/cs243/joeq/index.html

#### **Operators**

- □ Types of operators
  - Primitive operations: Moves, Adds, Bitwise AND, etc.
  - Memory access: Getfields and Getstatic
  - Control flow: Compares and conditional jumps, JSRs
  - Method invocation: OO and traditional
- □ Operators have suffixes indicating return type:
  - ADD\_I adds two integers
  - L, F, D, A, and V refer to longs, floats, doubles, references, and voids respectively
  - Operators may have \_DYNLINK (or %) appended, which means that a new class may need to be loaded

#### **Operands**

- □ Operands are split into 15 types
  - The ConstOperand classes (I, F, A, etc.) indicate constant values of the relevant type
  - RegisterOperands name pseudo-registers
  - MethodOperands and ParamListOperands are used to identify method targets
  - TypeOperands are passed to type-checking operators, or to "new" operators
  - TargetOperands indicate the target of a branch

#### Converting a Method to Quads

```
BB0 (ENTRY) (in: <none>, out: BB2)
BB2
       (in: BBO (ENTRY), out: BB3, BB4)
              TO int, R1 int, IConst: 10
   ADD I
2
   MOVE_I
              R3 int, T0 int
   ADD I
              TO int, R1 int, R3 int
4
   MOVE_I
              R6 int, T0 int
5
   IFCMP I
              R6 int, IConst: 2, LE, BB4
BB3
       (in: BB2, out: BB4)
   SUB_I
               TO int, R6 int, R3 int
7
   MOVE I
              R6 int, T0 int
       (in: BB2, BB3, out: BB1 (EXIT))
8
   RETURN I
              R6 int
       (EXIT) (in: BB4, out: <none>)
Exception handlers: []
Register factory: Local: (I=7, F=7, L=7, D=7, A=7)
                  Stack: (I=2, F=2, L=2, D=2, A=2)
```

#### Control Flow and CFGs

- □ joeq.Compiler.Quad.ControlFlowGraph encapsulates most of the information we need
  - Don't confuse with the ControlFlowGraph in joeq.Compiler.BytecodeAnalysis
- □ Generated from jq\_Methods by the underlying system's machinery

#### **Basic Blocks**

- □ Raw components of Control Flow Graphs
- □ Linked to predecessors and successors
- □ Contain a list of Quads
- And information about exception handlers
  - Which ones protect this basic block
  - Which blocks this one protects
- Exceptions violate traditional BB semantics
  - An exception can jump out of the middle of a basic block
  - We will ignore this subtlety

#### Safety Checks

- □ Java's safety checks are *implicit*: instructions may throw exceptions
- □ Joeq's safety checks are *explicit*: values of arguments are tested by operators such as NullCheck and BoundsCheck
  - Exceptions are thrown if checks fail
- □ When converting from bytecodes to quads, all necessary checks are automatically inserted

#### Iterating Over the Quads: QuadIterator

- Dealing with control flow graphs or basic blocks directly is tedious
- □ Dealing with individual quads tends to miss the forest for the trees
- □ Simple interface to iterate through all the quads in reverse post-order
- Predecessors and successors are still accessible
  jq\_Method m = ...
  ControlFlowGraph cfg = CodeCache.getCode(m);
  QuadIterator iter = new QuadIterator(cfg)
  while(iter.hasNext()) {
   Quad quad = (Quad)iter.next();
   if(quad.getOperator() instanceof Operator.Invoke) {
   processCall(cfg.getMethod(), quad);
   }
  }

# Developing a joeq Compiler Pass

### Writing and Running a Pass

- □ Passes themselves are written in Java, implementing various joeq interfaces
- □ Passes are invoked through library routines in the joeq.Main.Helper class

#### The joeq.Main.Helper Class

- □ joeq.Main.Helper provides a clean interface to the complexities of the joeq system
- □ load(String) takes the name of a class provides the corresponding jq\_Class
- □ runPass(target, pass) lets you apply any pass to a target that's at least that big

# Visitors in joeq

- □ joeq uses of the visitor design pattern
- ☐ The visitor for a level of the code hierarchy has methods visitFoo(code object) for each type of object in that level
- □ For some cases, types may overlap (e.g., visitStore and visitQuad) the methods will be called from least-general to most-general (i.e., visitStore before visitQuad)
- □ Visitor interfaces with more than one method have internal abstract classes called "EmptyVisitor" to simplify implementation

#### Visitors: Some Examples

```
public class QuadCounter extends QuadVisitor.EmptyVisitor
    {
      public int count = 0;
      public void visitQuad(Quad q) {
            count++;
      }
}

public class LoadStoreCounter extends
      QuadVisitor.EmptyVisitor {
      public int loadCount = 0, storeCount = 0;
      public void visitLoad(Quad q) { loadCount++; }
      public void visitStore(Quad q) { storeCount++; }
}
```

### Running a Pass

#### Summary

- □ We're using the joeq compiler system
- □ Review of Java VM's code hierarchy
- □ Review of joeq's code hierarchy
- QuadIterators
- □ joeq.Main.Helper
- □ Visitor pattern
- □ Defining and running passes

# Programming Assignment 1

- □ Implement a **basic data flow framework**
- □ We provide the interfaces for your framework
- □ Write the iterative Solver algorithm for any analysis matching these interfaces
- □ Phrase Reaching Definitions in terms that any Solver can understand

#### Programming Assignment 1

- □ Sample analyses are available in /usr/class/cs243/dataflow
- □ Flow.java contains the interfaces and the main() method
- □ ConstantProp.java contains classes that define a limited constant propagation algorithm
- ☐ Liveness.java contains classes that define a liveness detection algorithm

#### Programming Environment

- □ joeq.jar is provided
- □ We recommend you develop on Eclipse
- ☐ Your output must match ours on the publicly accessible Stanford Linux clusters
- □ Sample test case and matching output is provided
- □ The makefile will simplify your life

### The LivePC Engine

- □ Import a complete Eclipse development platform, including joeq (~800MB)
- ☐ Go to <a href="www.moka5.com">www.moka5.com</a> and download the LivePC Engine
- ☐ After installation, fetch the Java Program Analysis Toolset library
- □ (Your TAs are not proficient LivePC users)